Graph Reductions and Transformations

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Plan

- 1. Basic Graph Definitions
- 2. Important Theorems
- 3. Reductions
- 4. Transformations
- 5. What I'm Trying To Do

Technical Definition:

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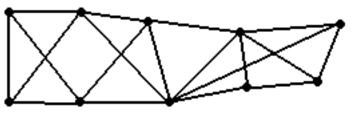
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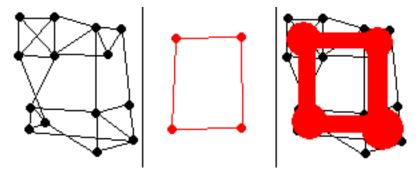
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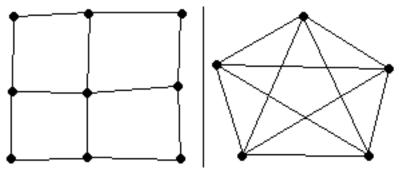
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▶ If you have a family F of graphs such that if a graph $G \in F$ then all minors of $G \in F$, then there exists a finite list of graphs L such that a graph H is in F iff H does not contain any graph of L as a minor.

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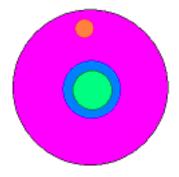
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Kuratowksi's Theorem:

▶ A graph is planar iff it does not have K_5 or $K_{3,3}$ as a minor.



 $Pink = all\ graphs$ $Green = our\ family$ $Blue = boundary\ set$ $Orange = graphs\ wearing\ pants$

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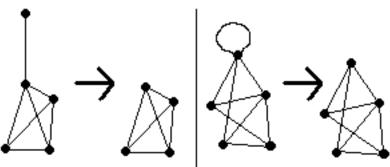
Useless Reductions

- ▶ If a vertex has only a single edge, we can delete that vertex and that edge.
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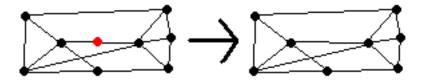


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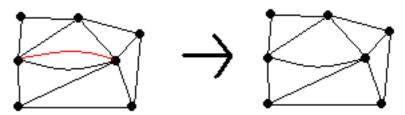


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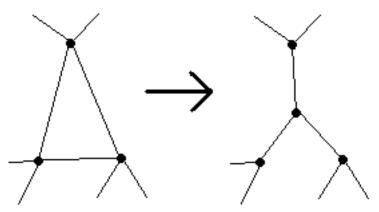


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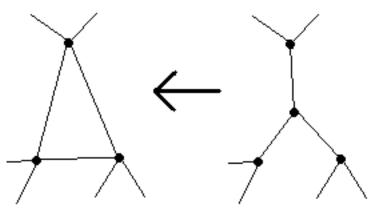


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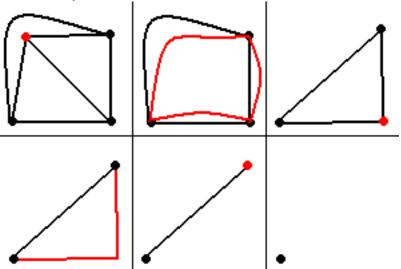
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- There must be a finite list of graphs that define the boundary between being reducible and not being reducible to isolated vertices.
- ➤ Yu (2004,2006) found 68 billion graphs in the boundary set. They fall into 20 families.

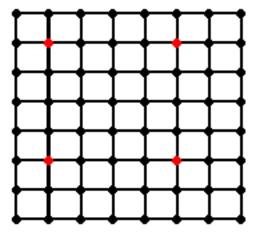
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- ▶ Doing this by considering how many neighbours the extra vertex has, and assuming that the planar graph is a grid.

The following graph can not be reduced to isolated vertices:

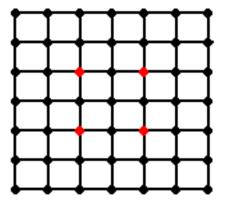
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The red vertices indicate where the extra vertex is attached to.

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- ▶ So far nothing new. Everything I have found that is not reducbile contains one of the graphs given by Yu as a minor.
- ▶ It is okay if I don't find more minors, and simply show that the existing list is complete.

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- ► There is a termination condition, since if I get to far away from planar, then the graph must not be reducible.
- Determining whether I've found a graph in Yu's list is not necessarily easy, so it would be nice to find a way to check this efficiently.

End Credits

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